

# *Security Mechanisms to support Real-Time Applications for OBAN*

Thomas J. Wilke



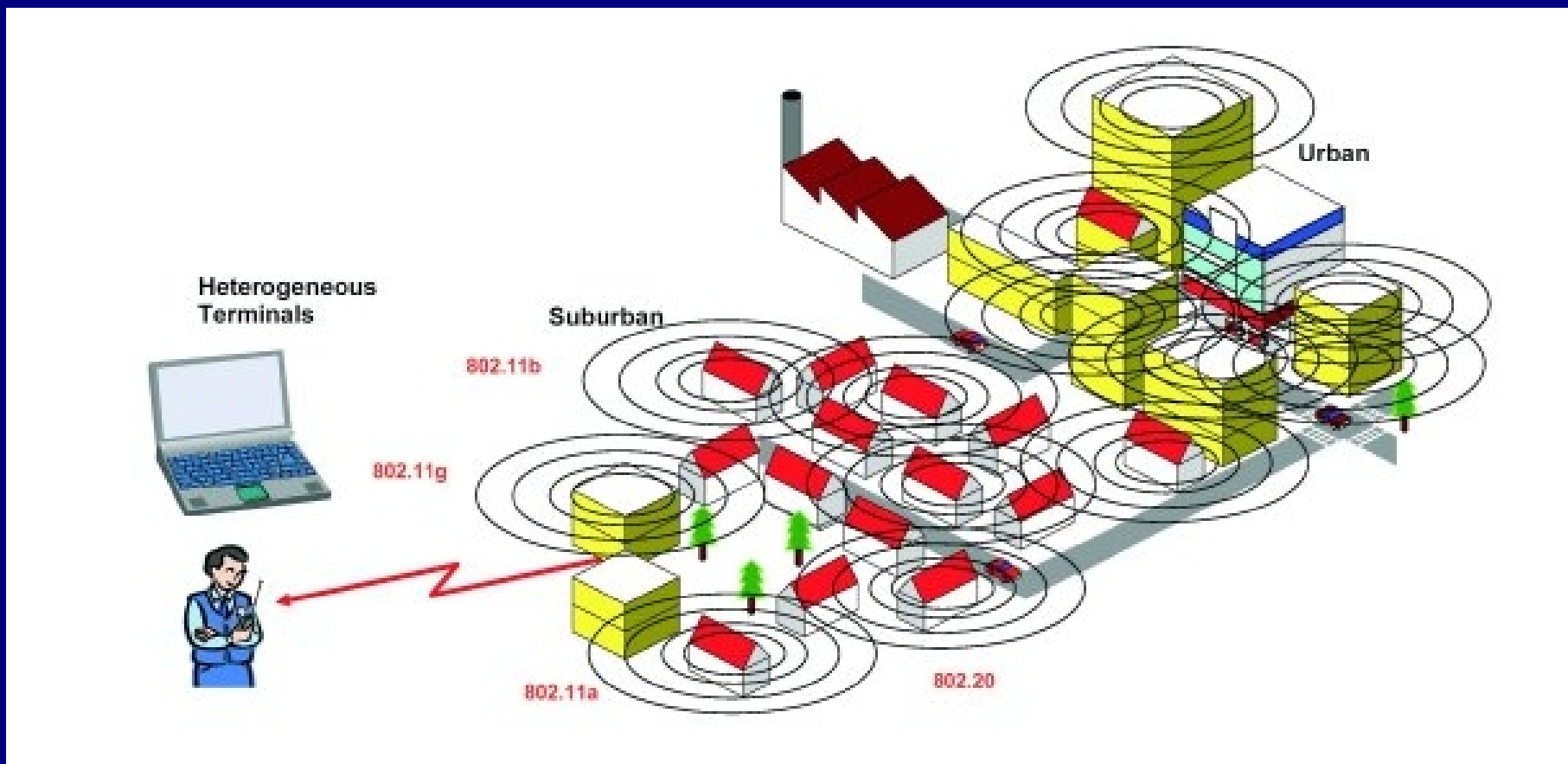
*Thomas J. Wilke Brussels, September 23rd*

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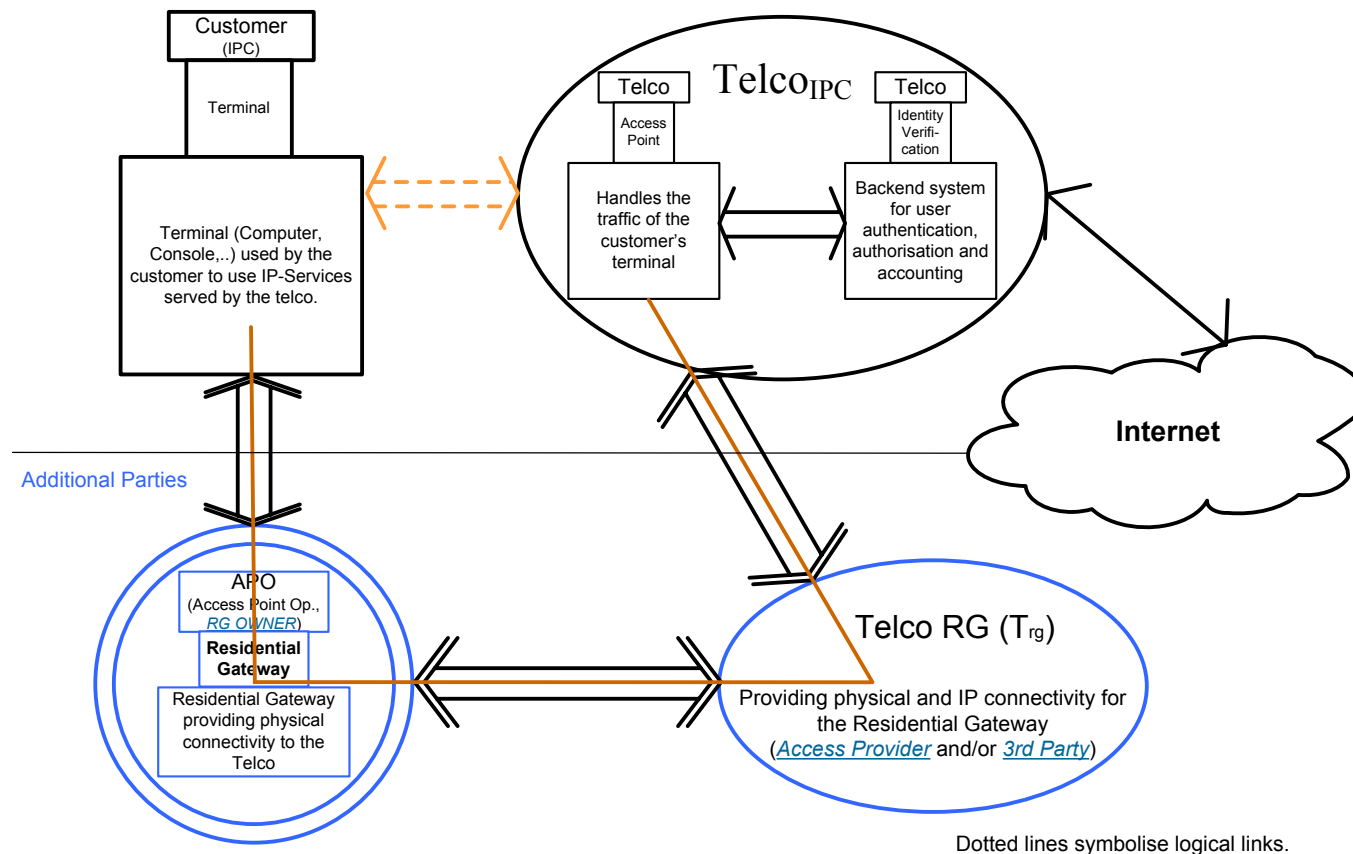
- OBAN Approach
- Security Goals
- Fast Handover and the Timing Problem
- Problem Solution: Time Shifted Security Computing
- Fast Handover Security Enforcement Mechanism
- Conclusion



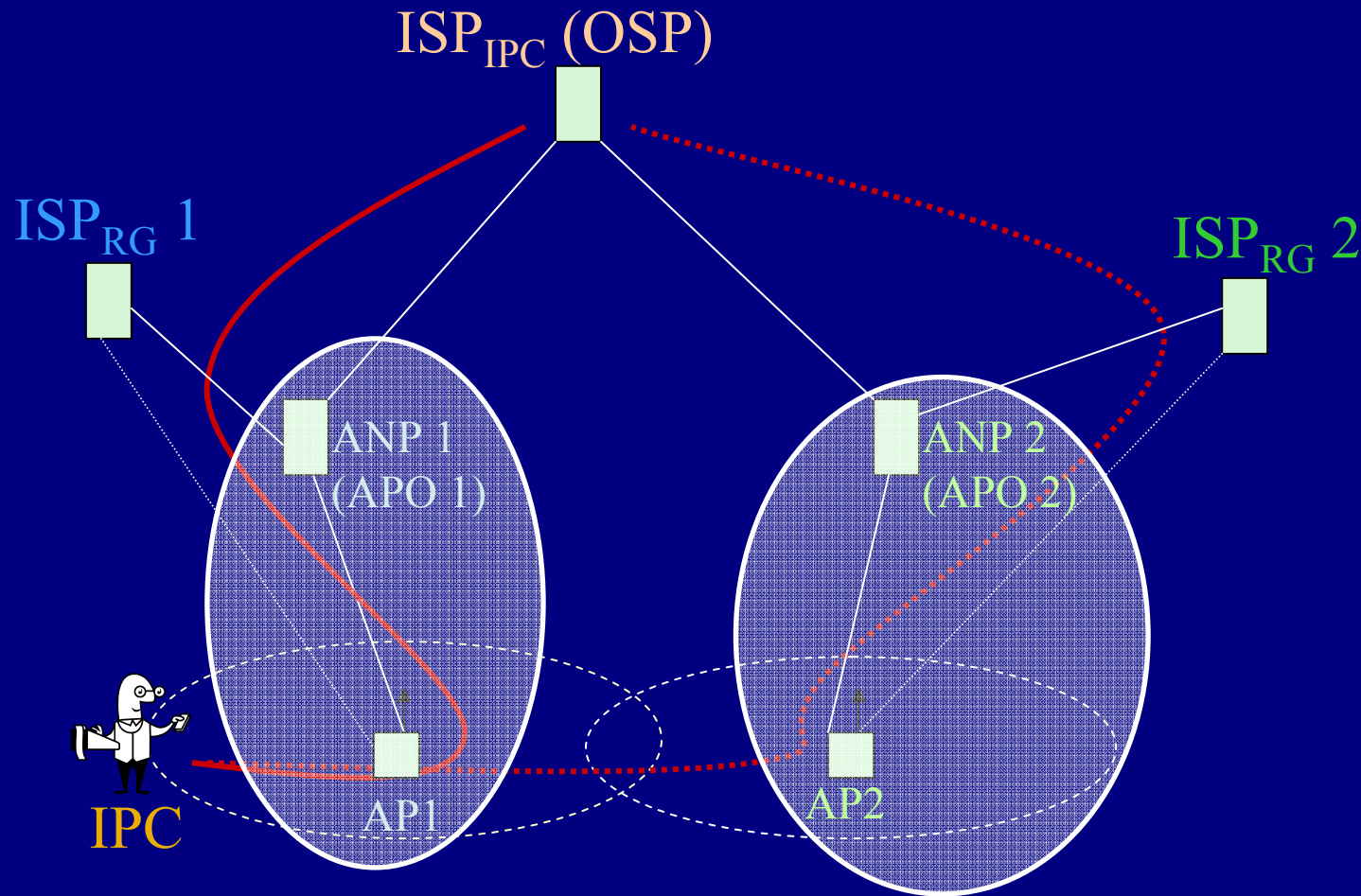
# The OBAN Approach



# OBAN Approach: Parties



# OBAN Parties: Mobility



# OBAN Security Goals

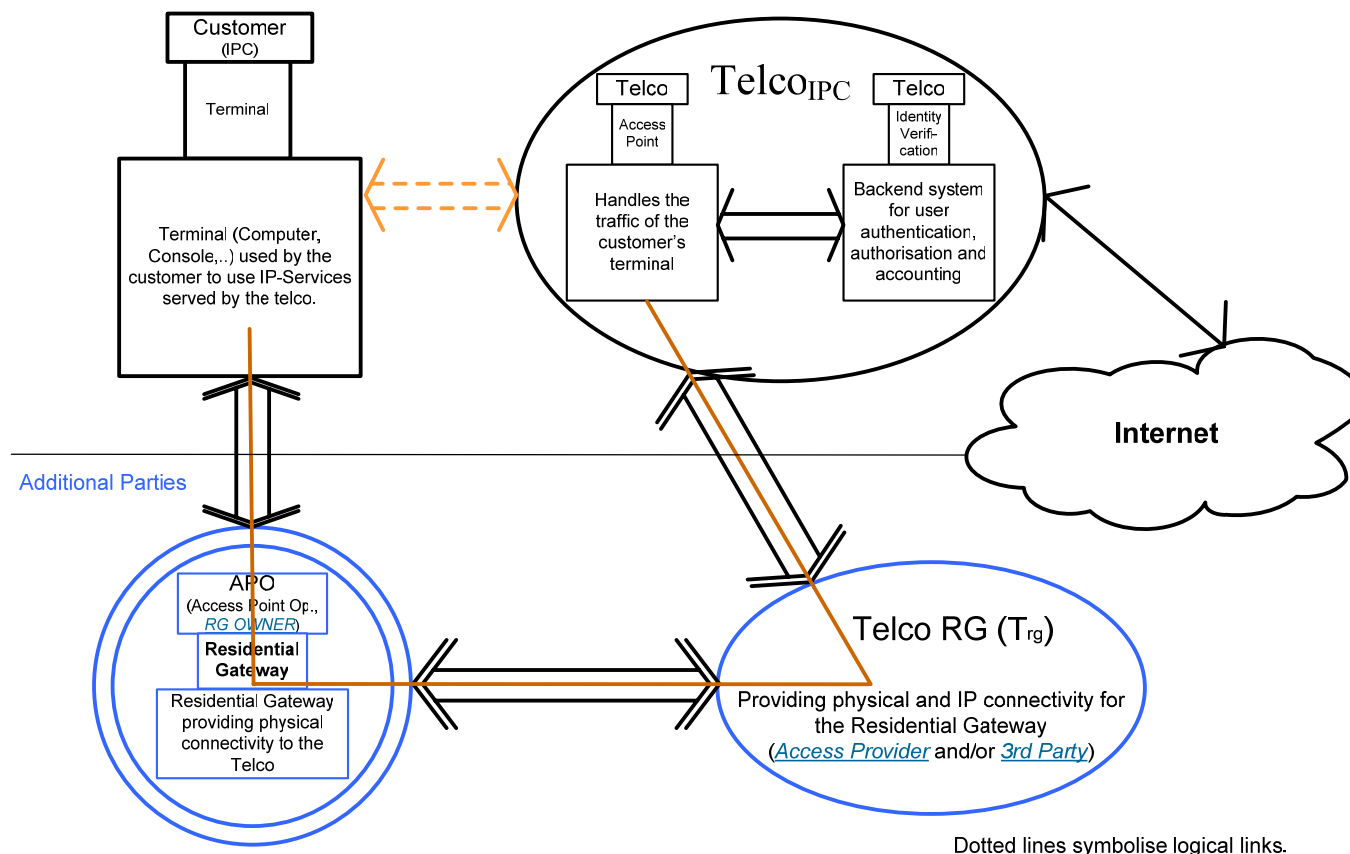
- Authentication
- Confidentiality
- Accountability
- Non-Repudiation
- Data Integrity
- Access Control
- Enhanced data protection and privacy
- Binding transparency

## Specific for Handover:

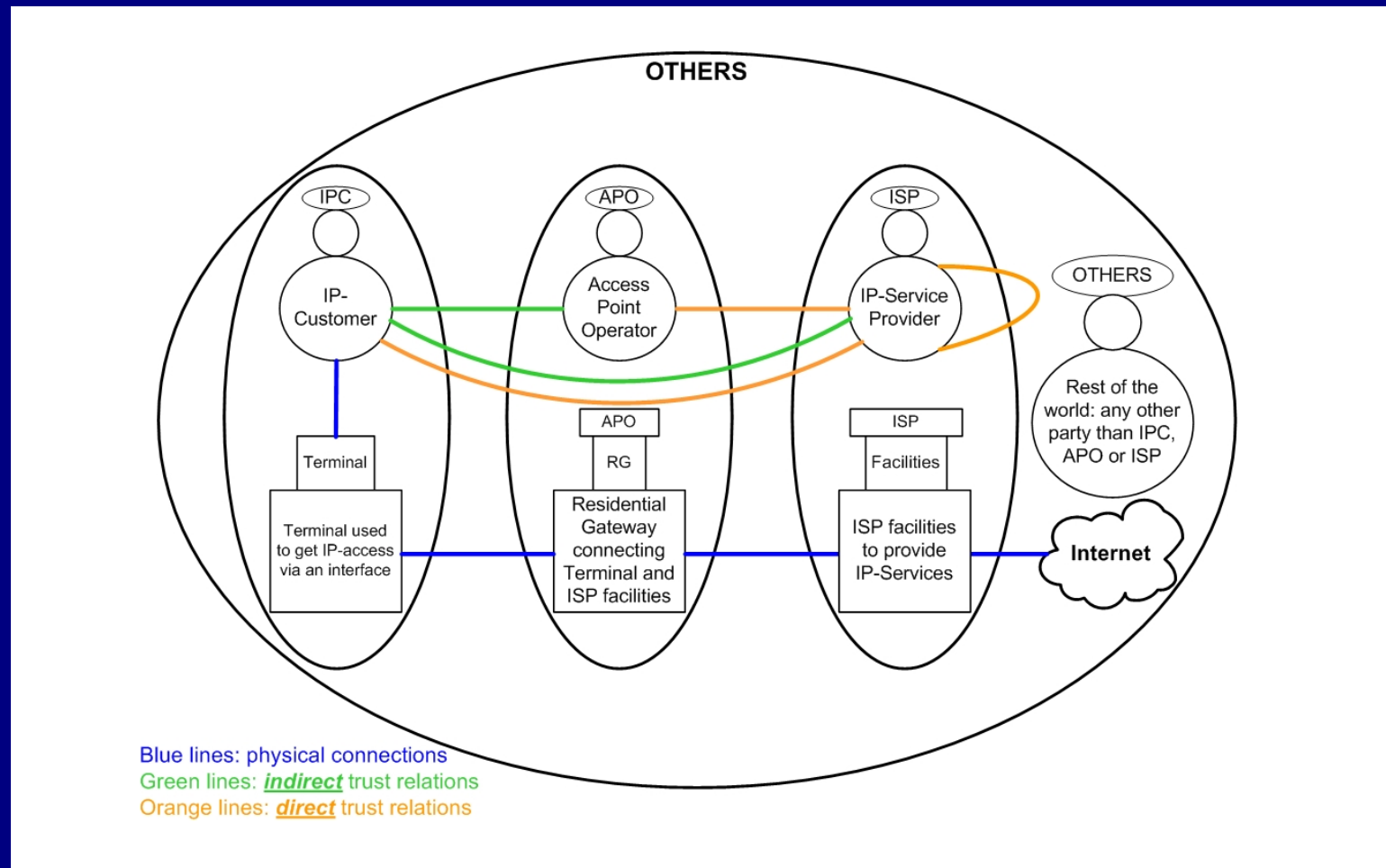
- Bridging Policy Enforcement



# OBAN Approach: Security Requirements

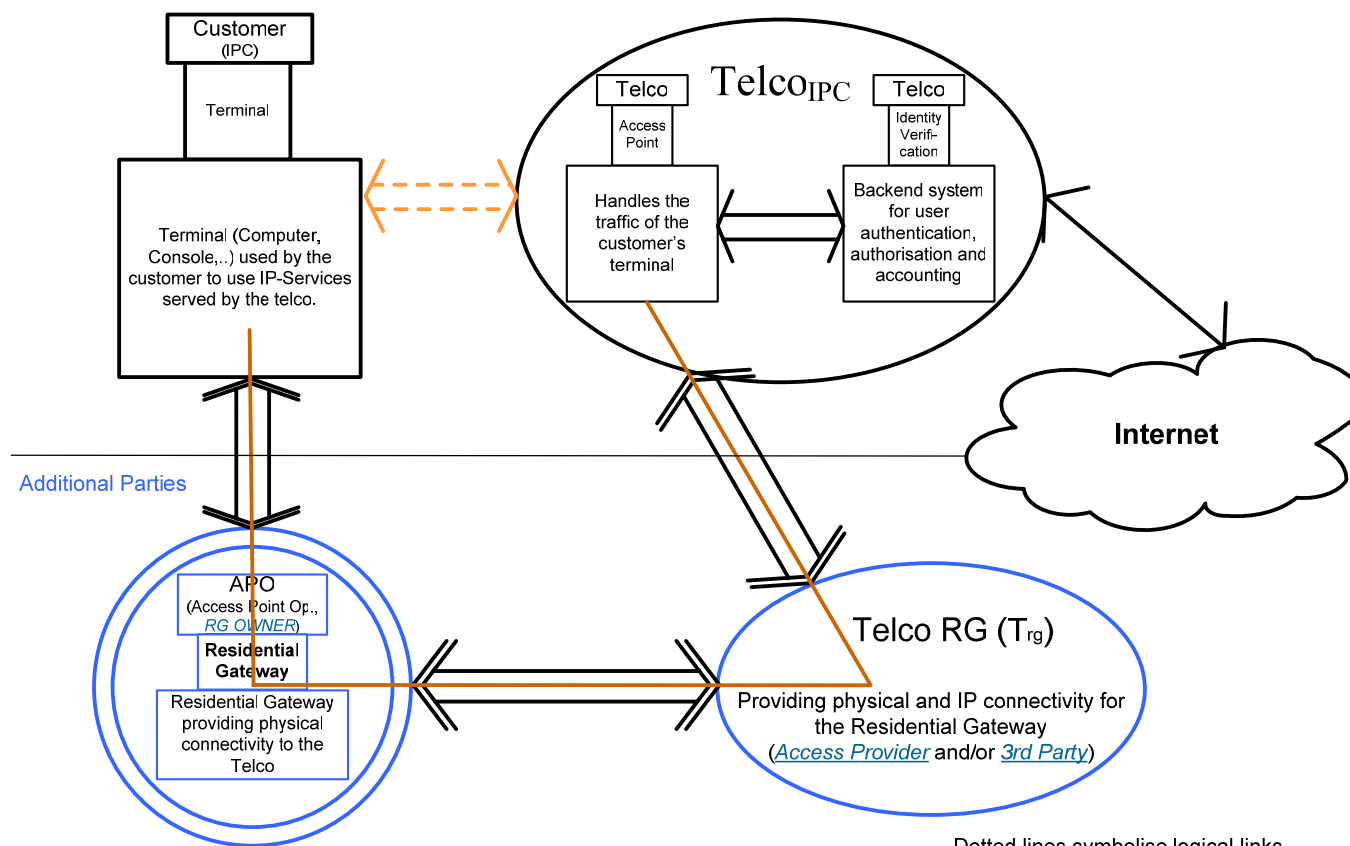


# OBAN Approach: Party Trust Relations

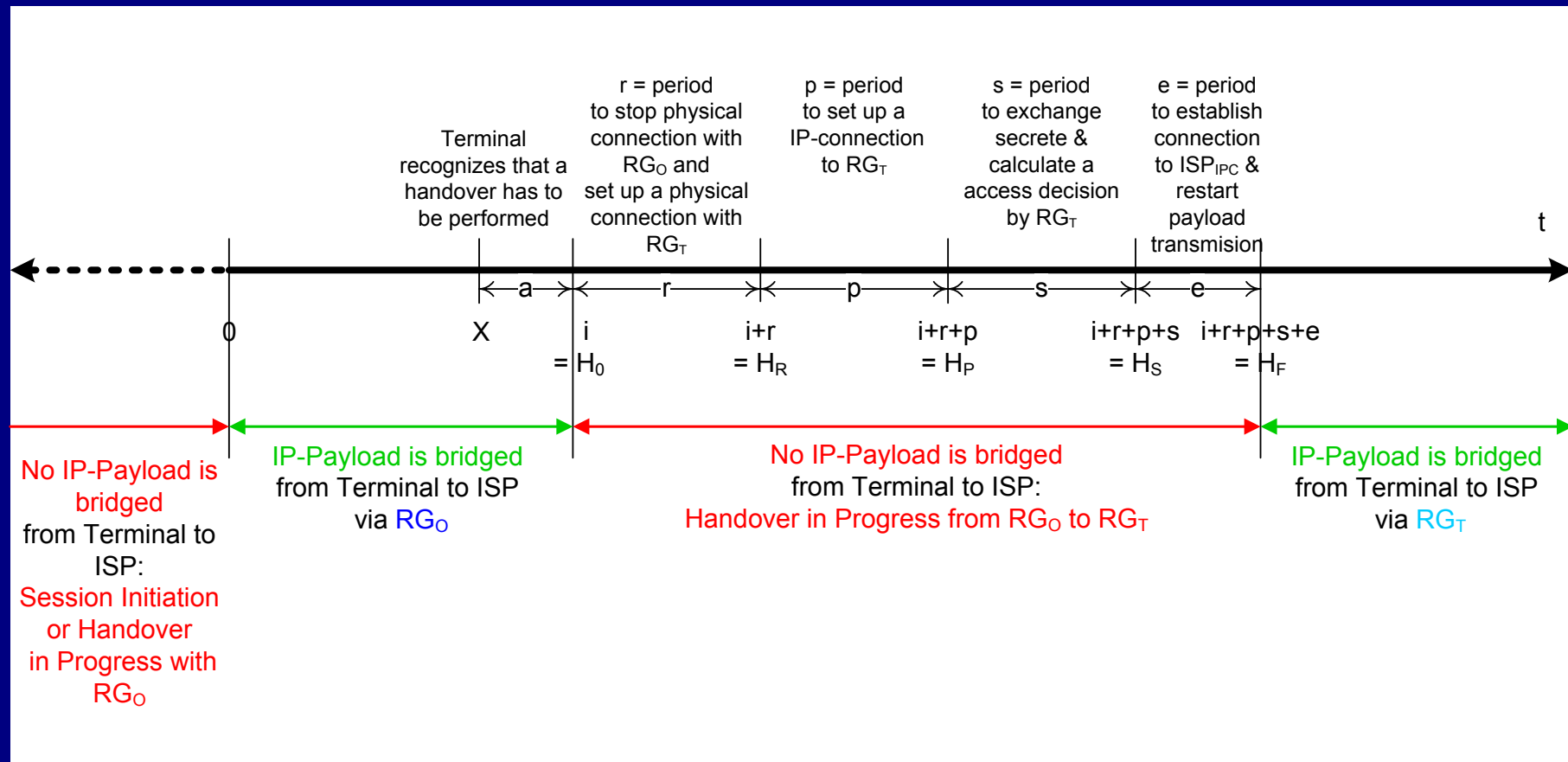




# OBAN Approach: Mutual Authentication



# Fast Handover and the Timing Problem



# Problem Solution: Time Shifted Security

## 2 Step divided security functionality processing:

- **Complex Security Functionality Processing:**  
Resource and time consuming are performed in parallel to other computations to prepare Real-Time Security Processing.
- **Real-Time Security Processing:**  
Very resource and time efficient processing performed based upon data computed ahead or after the Real-Time Processing event.



# Fast Handover Security Enforcement

## Mechanism: pre- & post-processing

Combination of 2 time shifted computing variants leading to 3 step computation mechanisms

- **1 step: Pre-Processing** (payload is bridged in parallel)
  - $\text{ISP}_{\text{IPC}}$ : Identifying candidate ISPs and RGs
  - $\text{ISP}_{\text{IPC}}$ : Terminal authorization
  - $\text{ISP}_{\text{RG}}$ : Authorization of candidate RGs
  - $\text{ISP}_{\text{RG}}$ : Terminal identity and authentication indication



# Fast Handover Security Enforcement

## Mechanism: pre- & post-processing

Combination of 2 time shifted computing variants  
leading to 3 step computation mechanisms

- **2 step: Real-Time-Processing** (no payload is bridged in parallel)
  - RG: Terminal identification
  - RG: Terminal access decision and enforcement



# Fast Handover Security Enforcement

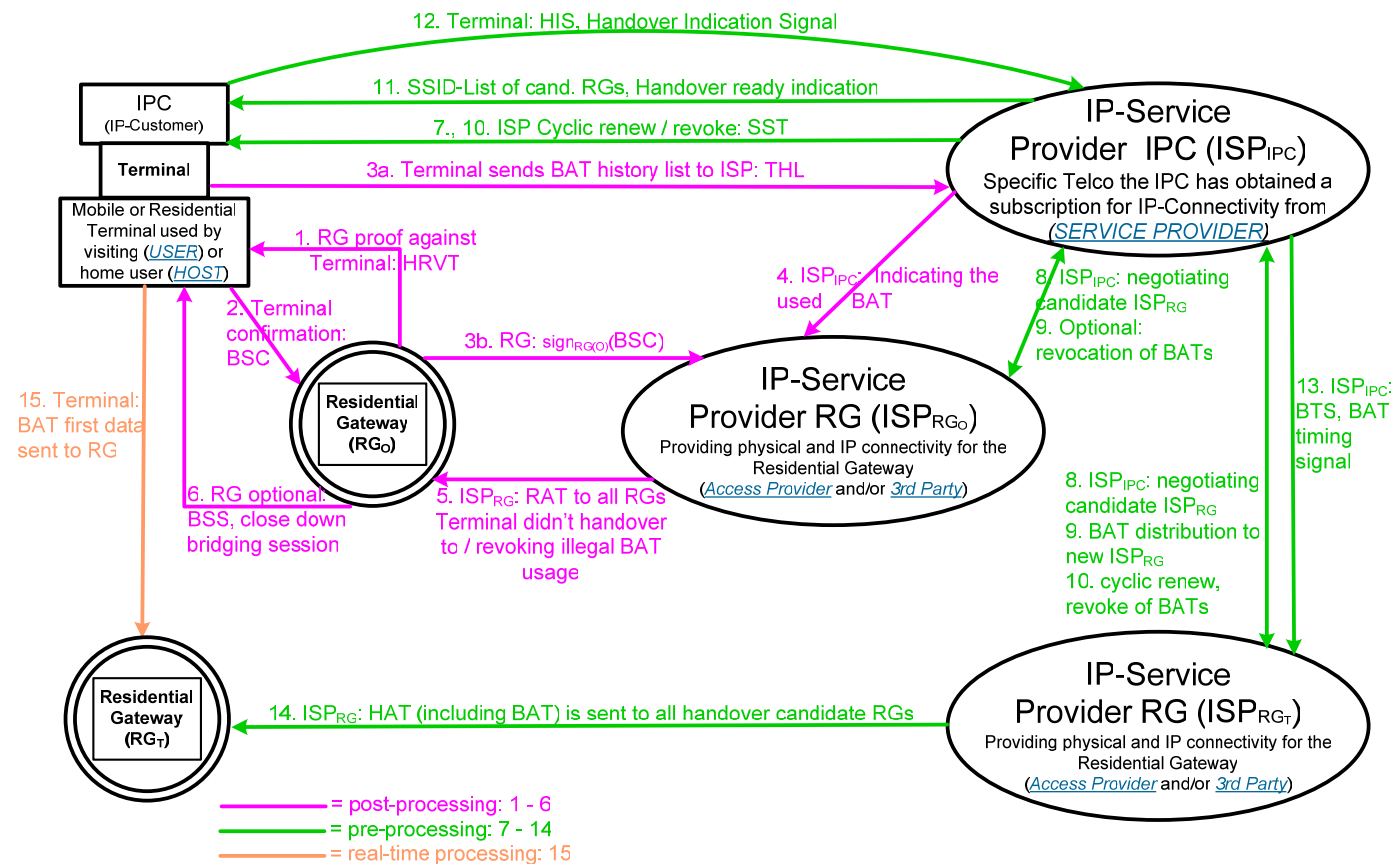
## Mechanism: pre- & post-processing

Combination of 2 time shifted computing variants leading to 3 step computation mechanisms

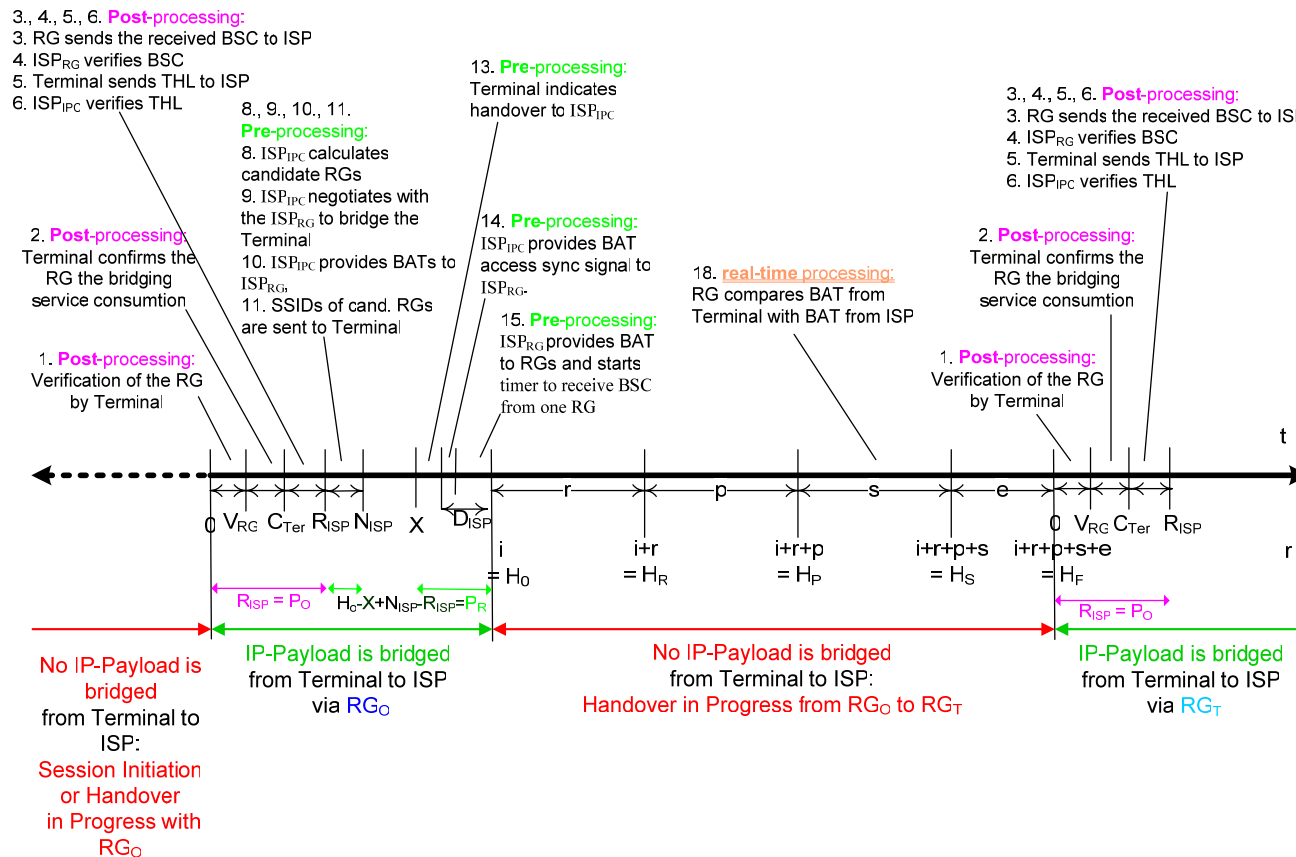
- **3 step: Post-Processing** (payload is bridged in parallel)
  - Terminal: RG is authenticated
  - RG: Bridging service consumption confirmation from Terminal
  - $ISP_{IPC/RG}$  : Service consumption legitimacy verification



# Fast Handover Security Enforcement Mechanism: party perspective

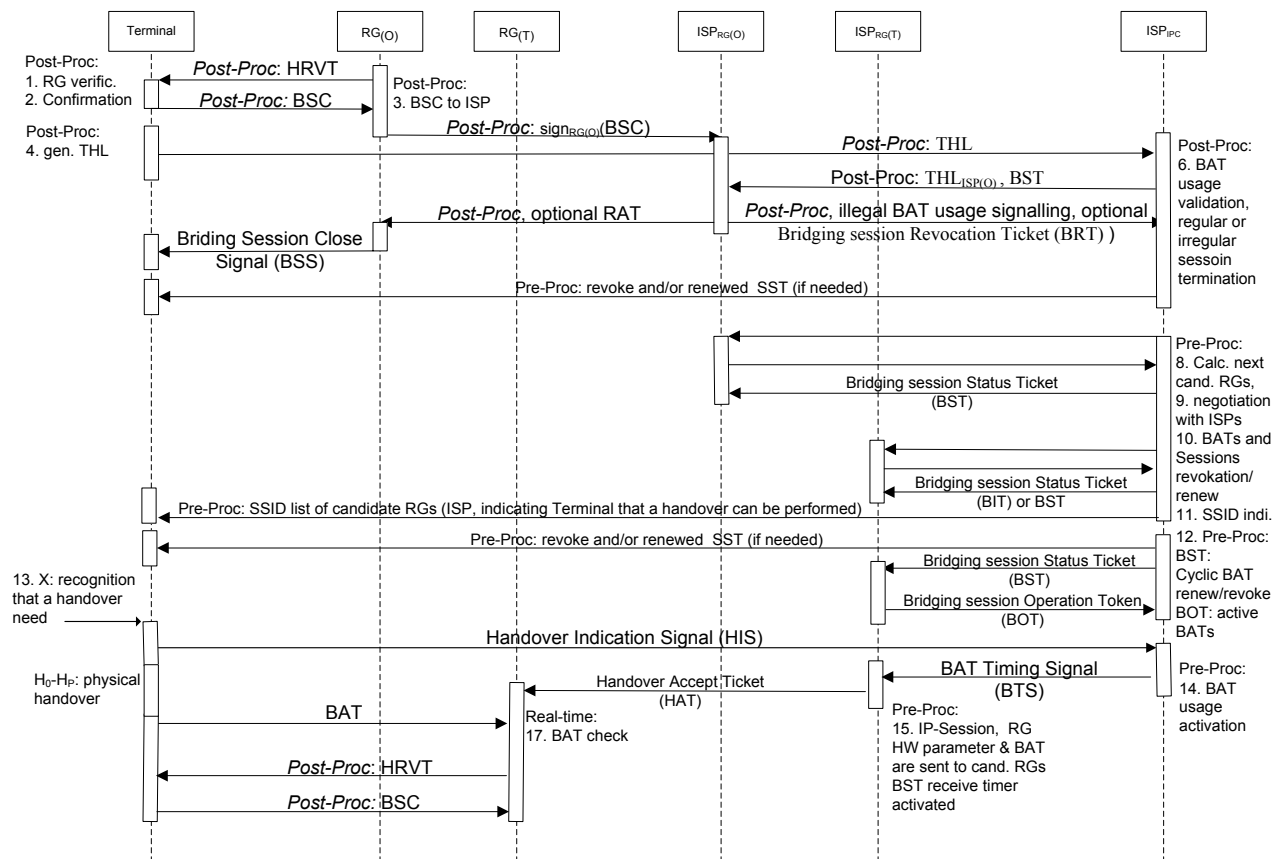


# Fast Handover Security Enforcement Mechanism: time-scale perspective





# Fast Handover Security Enforcement Mechanism



# Characteristics of the Mechanisms

The presented mechanisms can be characterized by the following items:

- A Bridging Access Ticket (BAT) describes a handover in respect of the following aspects:
  - BAT is created and distributed by  $ISP_{IPC}$
  - BAT is only valid for a definite Handover within a limited period,
  - BAT is distributed to 1 or a max.  $n$  RGs,
  - BAT can only be used by one Terminal
- Parties controlling the access of a Terminal (in hierarchical order)
  - $ISP_{IPC}$
  - $ISP_{RG}$
  - RG



# Conclusion

A mechanism has been presented enabling a fast handover solution for OBAN which

- Holds all security goals demanded by the business models (WP1)
- Holds all security requirements defined within the basic OBAN security architecture
- Enables in addition a bridging policy enforcement to ensure a fair and economic resource usage of RGs and ISPs



# References

- OBAN Security Architecture, document D10.1
- Document file:  
OBAN-WP1-SCOM-16c\_v03-tn2.ppt
- Document file:  
OBAN-WP2-TUB-311d-DIS.pdf



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